Substrate Storage

Deep dive.

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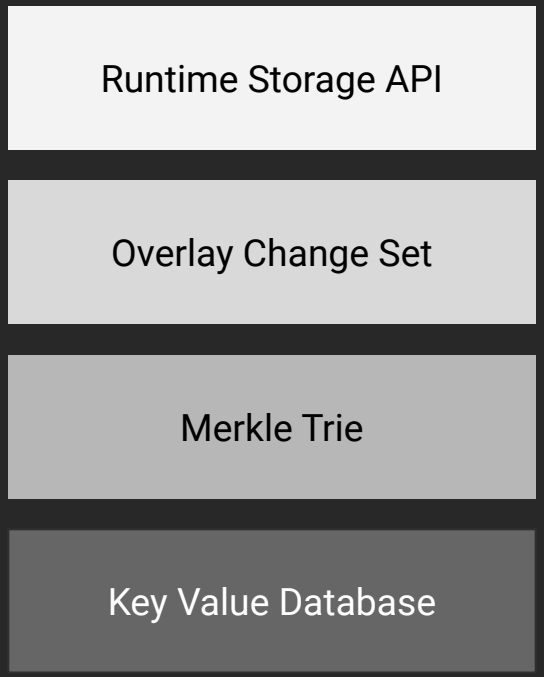
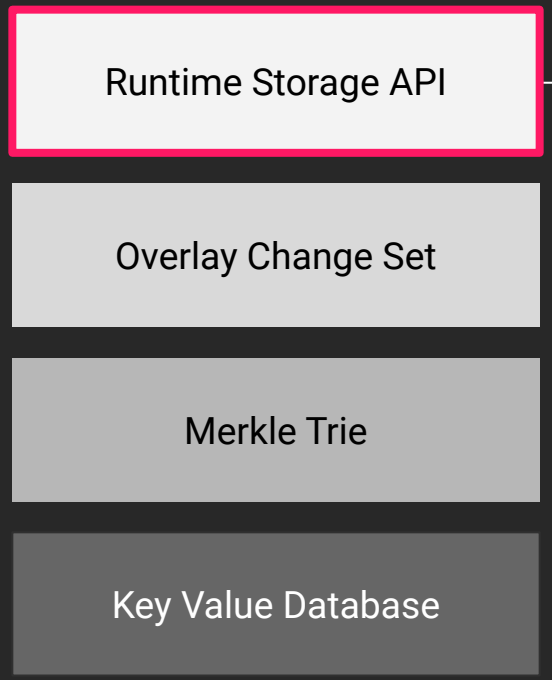
Abstractions of Substrate Storage



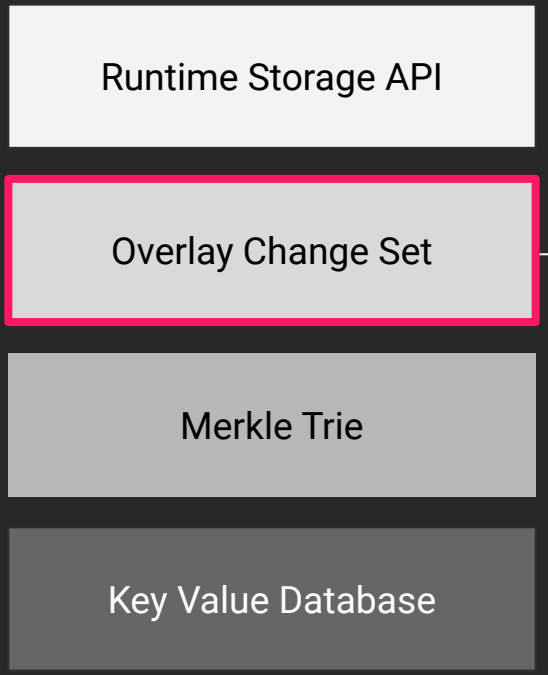
High Level Overview

Abstractions of Substrate Storage



- sp-io can write to storage with a given key + value
- Easy APIs generated through decl_storage! macro
- StorageValue, StorageMap, StorageDoubleMap, etc...

Abstractions of Substrate Storage



- Stages changes to the underlying DB.
- Overlay changes are committed once per block.
- Two kinds of changes:
  - Prospective Changes - what may happen.
  - Committed Changes - what will happen.

Abstractions of Substrate Storage

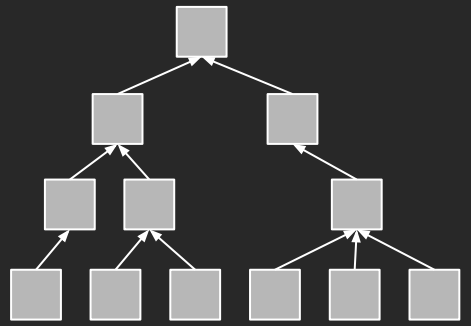
Runtime Storage API

Overlay Change Set

Merkle Trie

Key Value Database

- a.k.a. HashDB
- paritytech/trie
- Data structure on top of KVDB
- Arbitrary Key and Value length
- Nodes are Branches or Leaves



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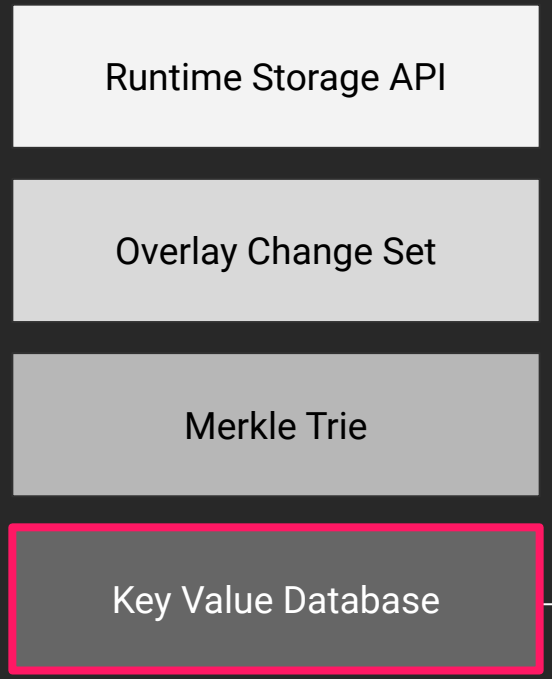
Abstractions of Substrate Storage



- a.k.a. KVDB
- Implemented with RocksDB
- Hash -> Vec<u8>
- Substrate: Blake2 256

Key (Hash 256)	Value (Vec<u8>)
0x0fd923ca5e7...	[00]
0x92cdf578c47...	[01]
0x31237cdb79...	[02]
0x581348337b...	[03]

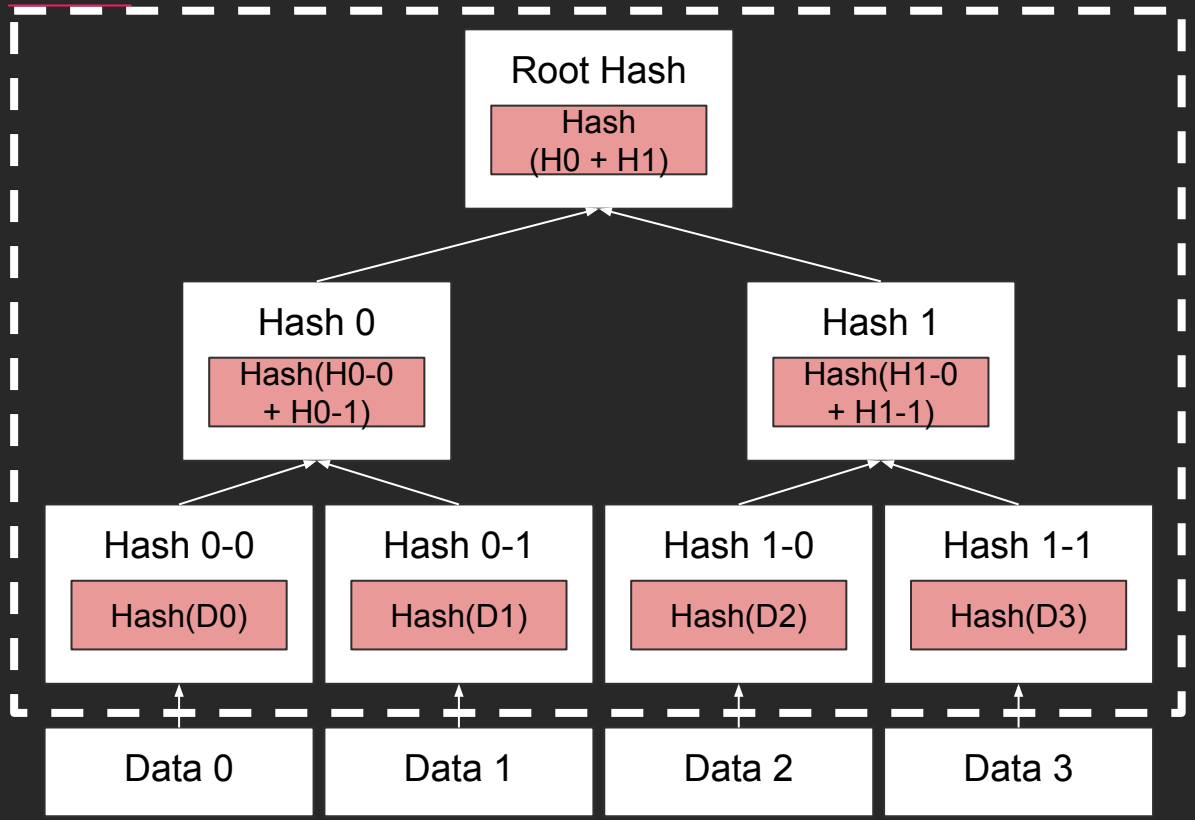
Two Kinds of Keys!

- Trie key path
- KVDB key hash

Don’t worry, we will come back to this...

Substrate uses a
Base-16 Patricia Merkle Trie

Merkle Tree



Root Node

Can be used to verify two trees are the same.

Branch Nodes

Leaf Nodes

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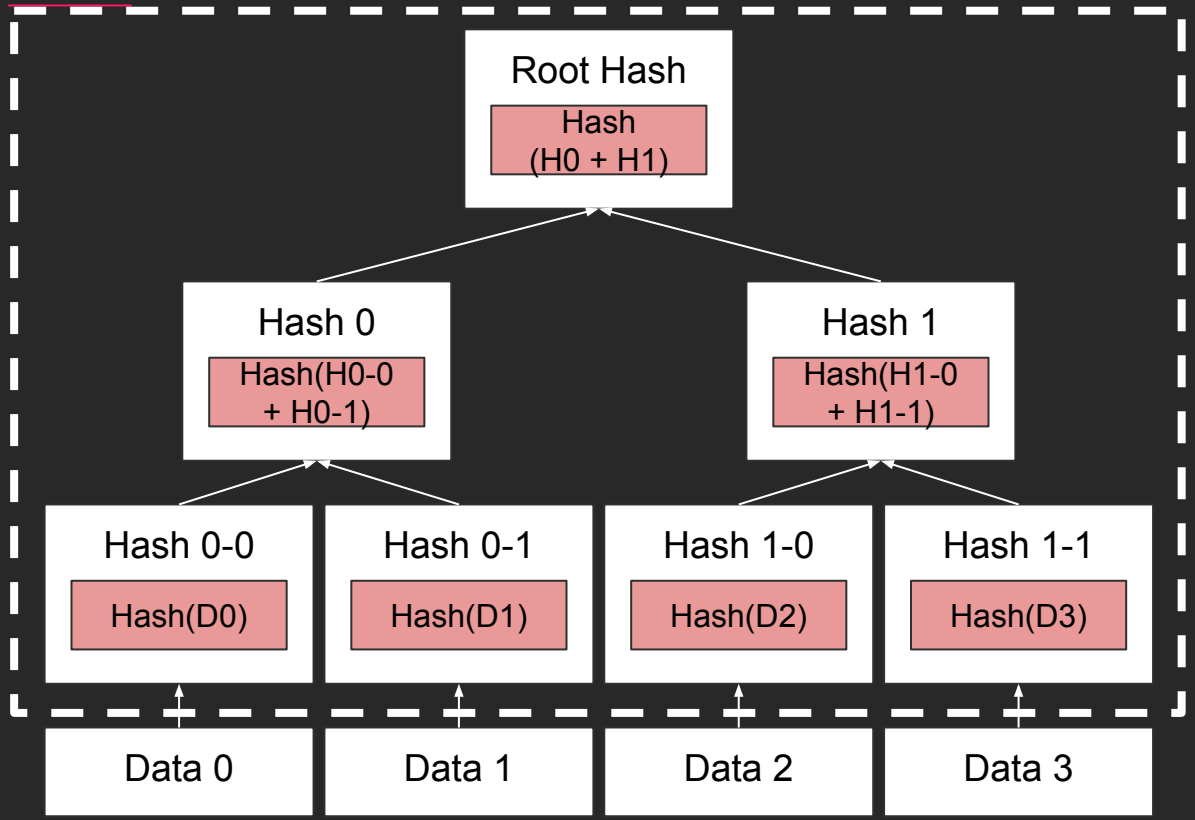
Merkle Tree



Merkle tree allows you to more easily prove that some data exists within the tree with a “Merkle Proof”.

More about that later.

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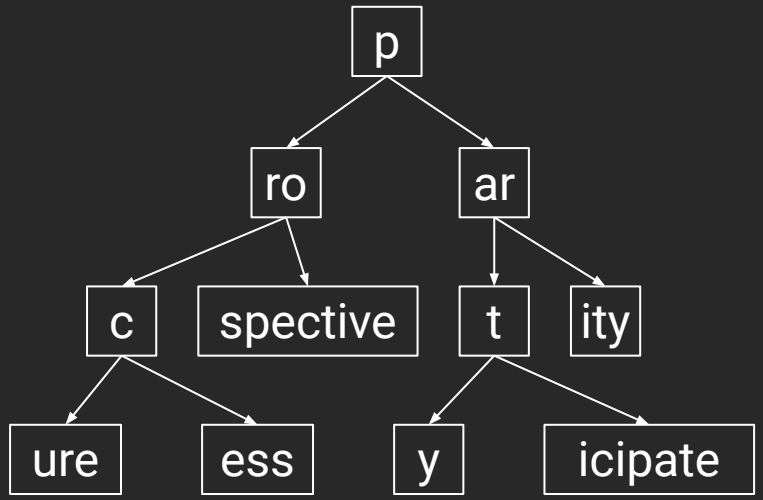
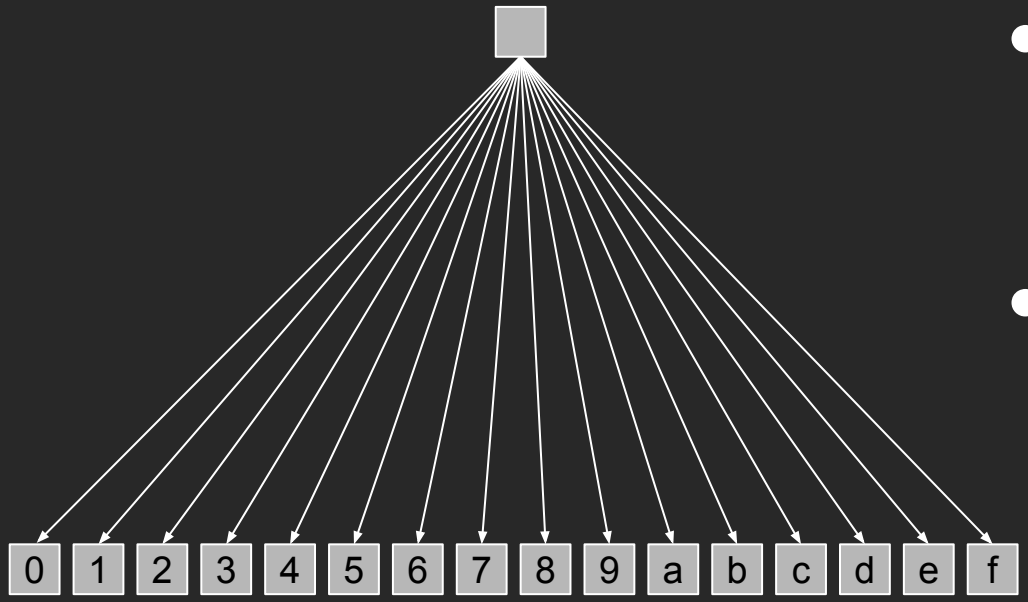
Patricia Trie



1. parity	2. participate	3. party
4. process	5. procure	6. prospective

- Position in the tree defines the associated key.
- Space optimized for elements which share a prefix.

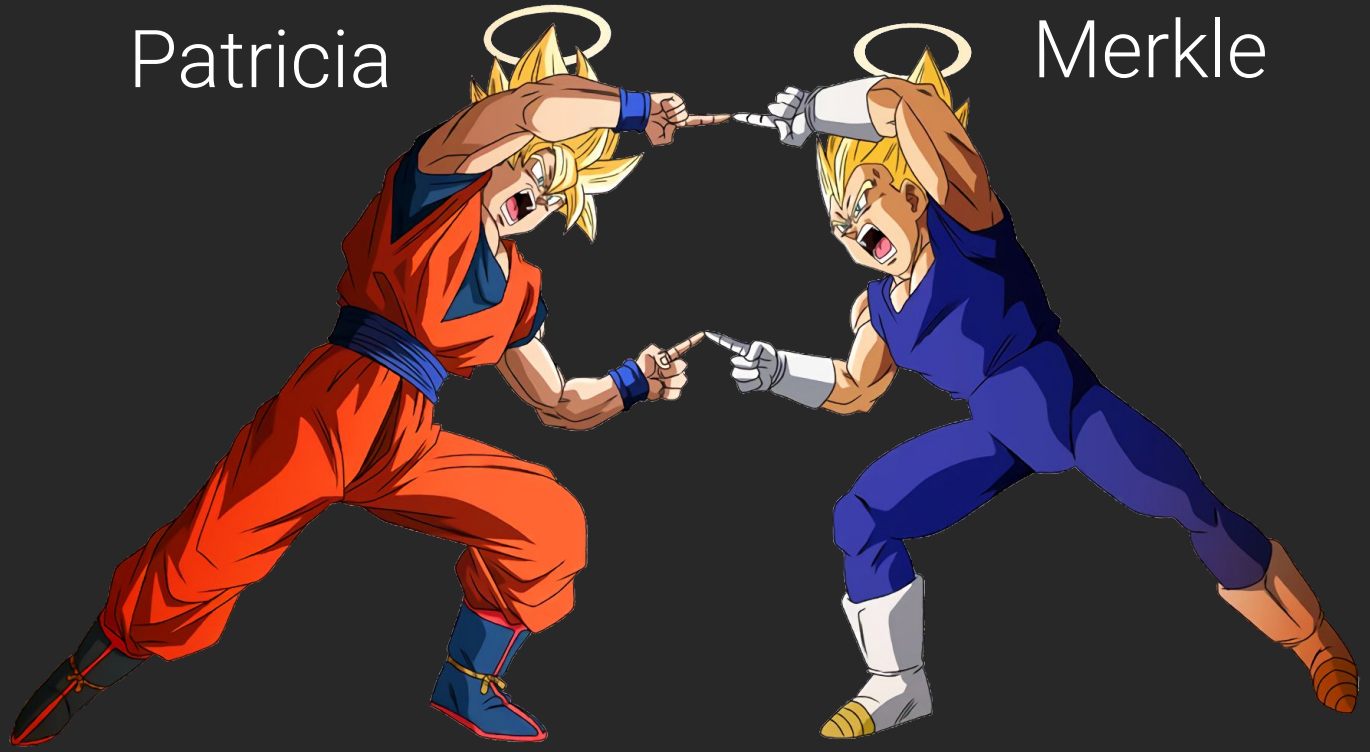
Beyond Binary Trees



- Branches can have more than two children.
- Everything is the same, just scaled up.

A single hex character is called a “nibble”.

Patricia



Merkle

Creation of the Patricia Merkle Trie

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Let’s get visual.

What we will be working with...

Literal KVDB Table

Key	Value
0x8f35a27d9...	[BRANCH]
0x2ebcd78e8...	[LEAF 00]
0x27434bcd0...	[BRANCH w/ VAL 01]
0x802c9c18c...	[LEAF 02]
0x986d278c5...	[LEAF 03]

Types of Nodes

Prefix	Type
00	Empty
01	Leaf
10	Branch w/o value
11	Branch w value

Virtual Trie Table

Trie Key Path							Value
a	7						[BRANCH]
a	7	1	1	3	5	5	[LEAF 00]
a	7	7	d	3			[BRANCH w/ VAL 01]
a	7	7	d	3	3	7	[LEAF 02]
a	7	7	d	3	9	7	[LEAF 03]

Node Structure

Trie Node			
header	key	children	value

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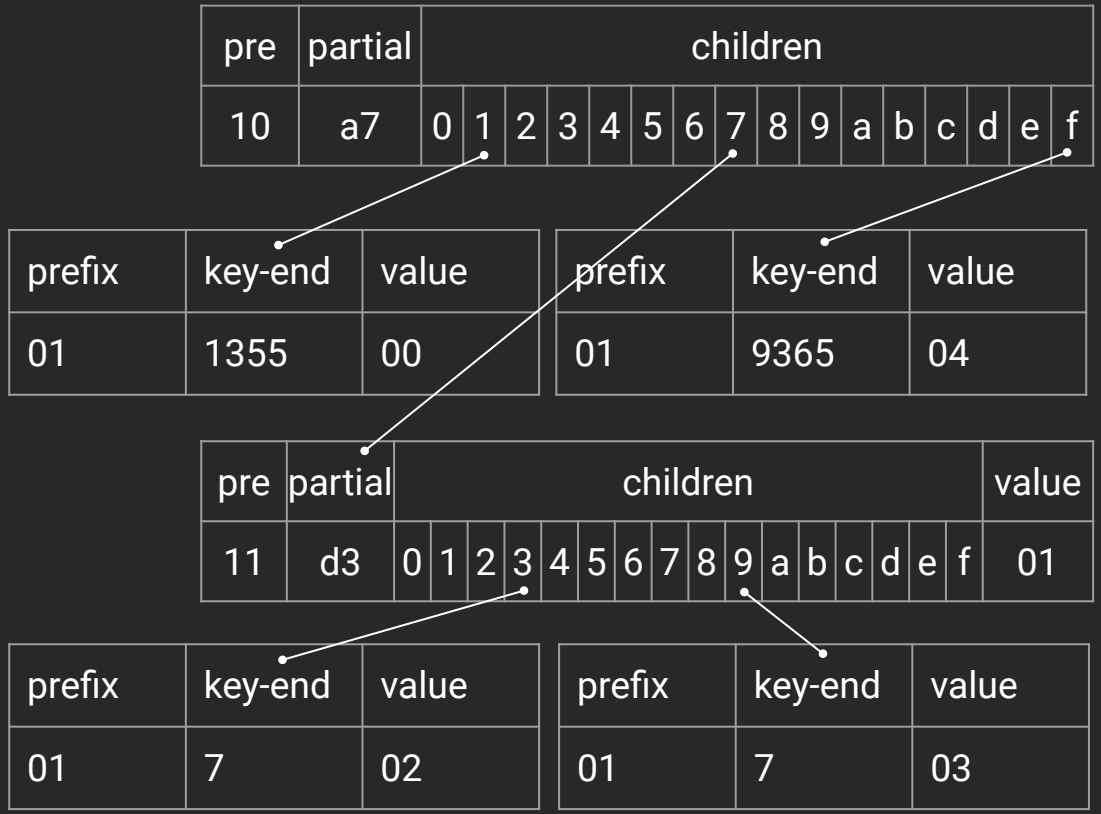
Visual of the Substrate State Trie



Trie Key Path							Value
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a	7	1	1	3	5	5	[LEAF 00]
a	7	7	d	3			[BRANCH w/ VAL 01]
a	7	7	d	3	3	7	[LEAF 02]
a	7	7	d	3	9	7	[LEAF 03]
a	7	f	9	3	6	5	[LEAF 04]

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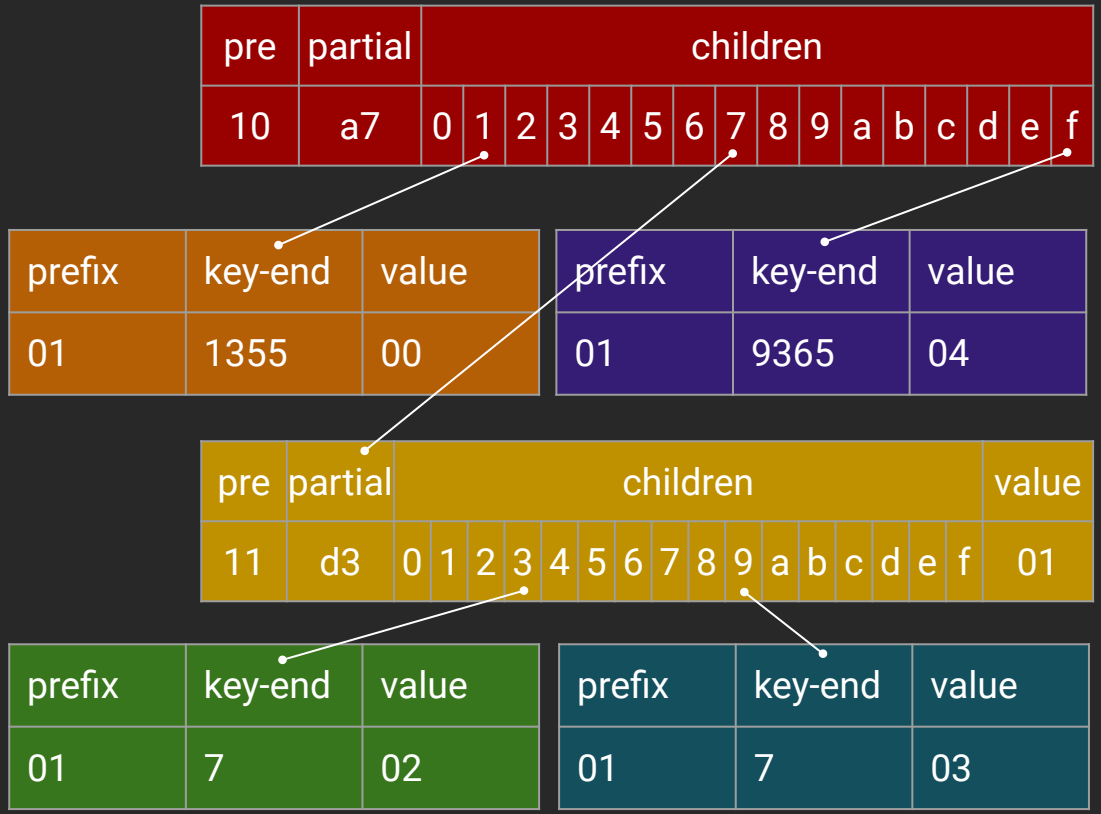
Visual of the Substrate State Trie



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a	7	1	1	3	5	5	[LEAF 00]
a	7	7	d	3			[BRANCH w/ VAL 01]
a	7	7	d	3	3	7	[LEAF 02]
a	7	7	d	3	9	7	[LEAF 03]
a	7	f	9	3	6	5	[LEAF 04]

All nodes are present.

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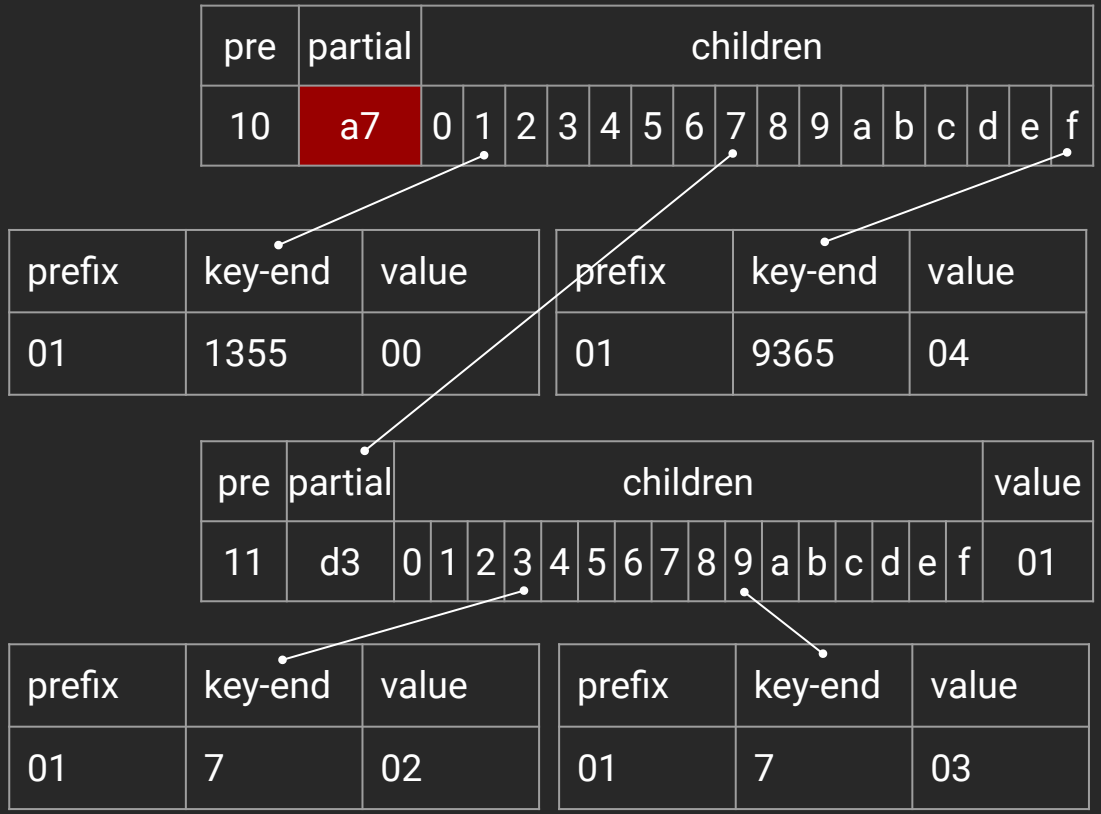
Visual of the Substrate State Trie



Trie Key Path							Value
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a	7	1	1	3	5	5	[LEAF 00]
a	7	7	d	3			[BRANCH w/ VAL 01]
a	7	7	d	3	3	7	[LEAF 02]
a	7	7	d	3	9	7	[LEAF 03]
a	7	f	9	3	6	5	[LEAF 04]

Nodes with a shared path are children of a branch.

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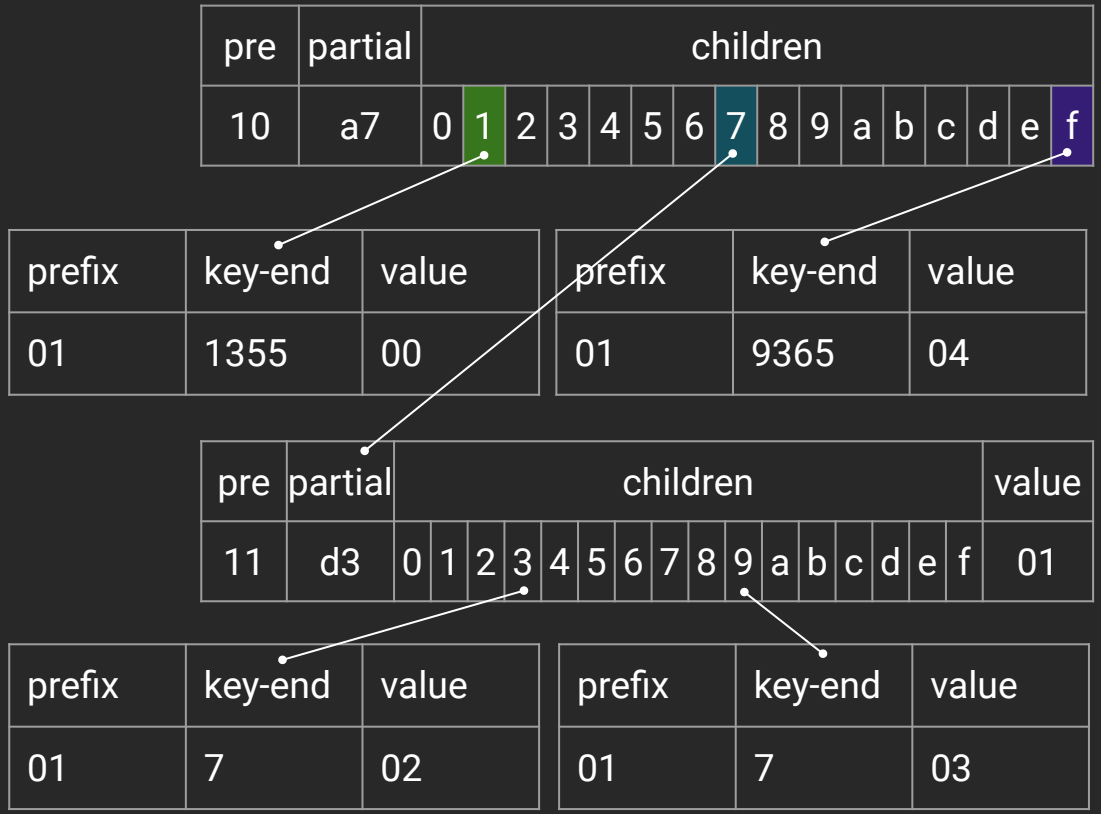
Visual of the Substrate State Trie



Trie Key Path							Value
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a	7	7	d	3			[BRANCH w/ VAL 01]
a	7	7	d	3	3	7	[LEAF 02]
a	7	7	d	3	9	7	[LEAF 03]
a	7	f	9	3	6	5	[LEAF 04]

You can then progress by looking at the children of the branch.

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Visual of the Substrate State Trie

pre	partial	children														
10	a7	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f

prefix	key-end	value
01	1355	00

prefix	key-end	value
01	9365	04

pre	partial	children																value
11	d3	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f	01

prefix	key-end	value
01	7	02

prefix	key-end	value
01	7	03

Trie Key Path							Value
a	7						[BRANCH]
a	7	1	1	3	5	5	[LEAF 00]
a	7	7	d	3			[BRANCH w/ VAL 01]
a	7	7	d	3	3	7	[LEAF 02]
a	7	7	d	3	9	7	[LEAF 03]
a	7	f	9	3	6	5	[LEAF 04]

This is a KVDB look up!

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Visual of the Substrate State Trie

pre	partial	children														
10	a7	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f

prefix	key-end	value
01	1355	00

prefix	key-end	value
01	9365	04

pre	partial	children																value
11	d3	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f	01

prefix	key-end	value
01	7	02

prefix	key-end	value
01	7	03

Trie Key Path							Value
a	7						[BRANCH]
a	7	1	1	3	5	5	[LEAF 00]
a	7	7	d	3			[BRANCH w/ VAL 01]
a	7	7	d	3	3	7	[LEAF 02]
a	7	7	d	3	9	7	[LEAF 03]
a	7	f	9	3	6	5	[LEAF 04]

You can have a branch which also contains a value!

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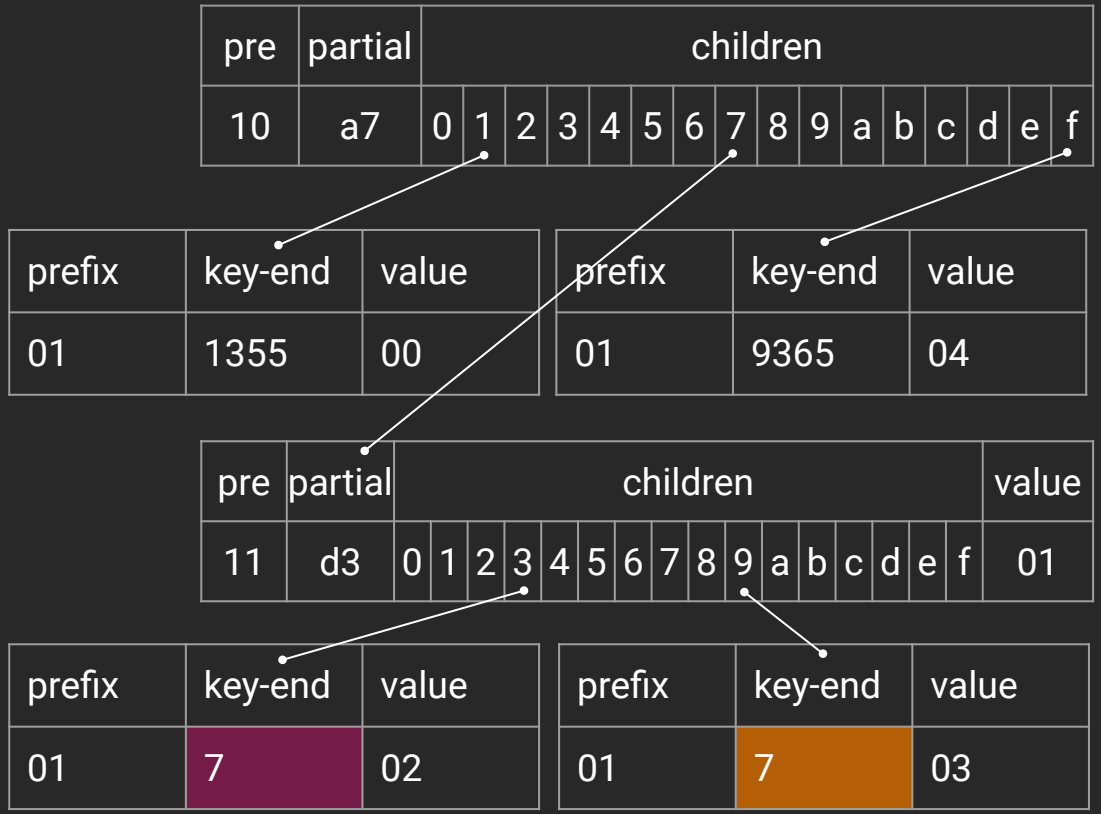
Visual of the Substrate State Trie



Trie Key Path							Value
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a	7	1	1	3	5	5	[LEAF 00]
a	7	7	d	3			[BRANCH w/ VAL 01]
a	7	7	d	3	3	7	[LEAF 02]
a	7	7	d	3	9	7	[LEAF 03]
a	7	f	9	3	6	5	[LEAF 04]

You reach the end when there are no more branches.

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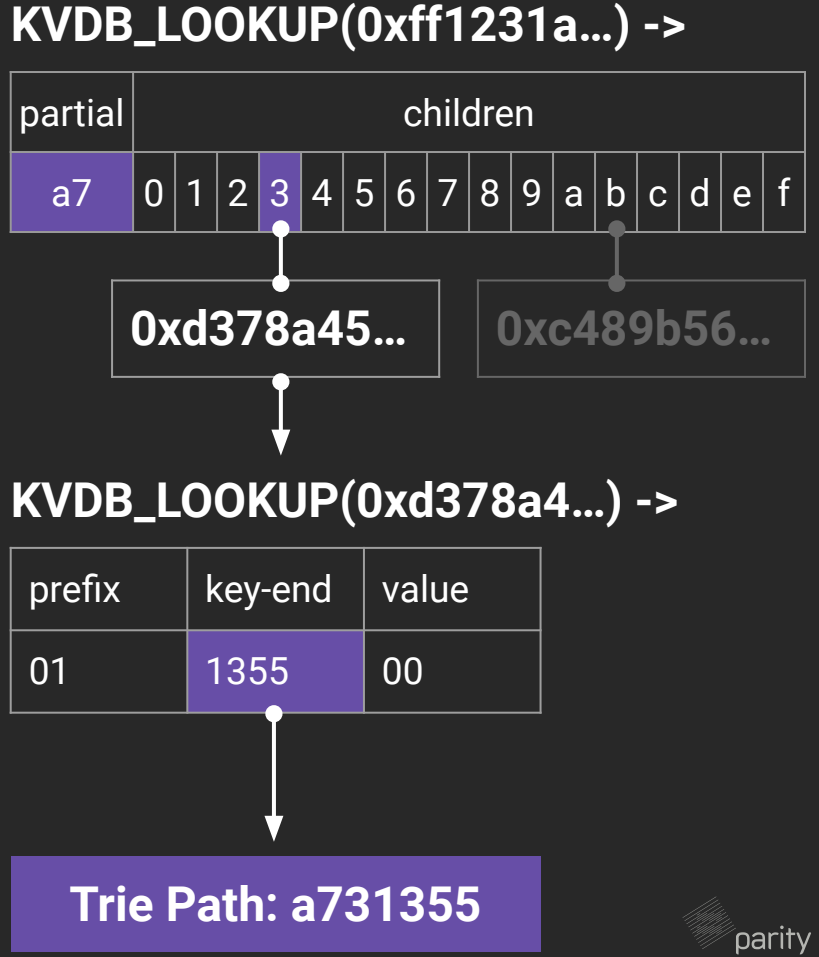
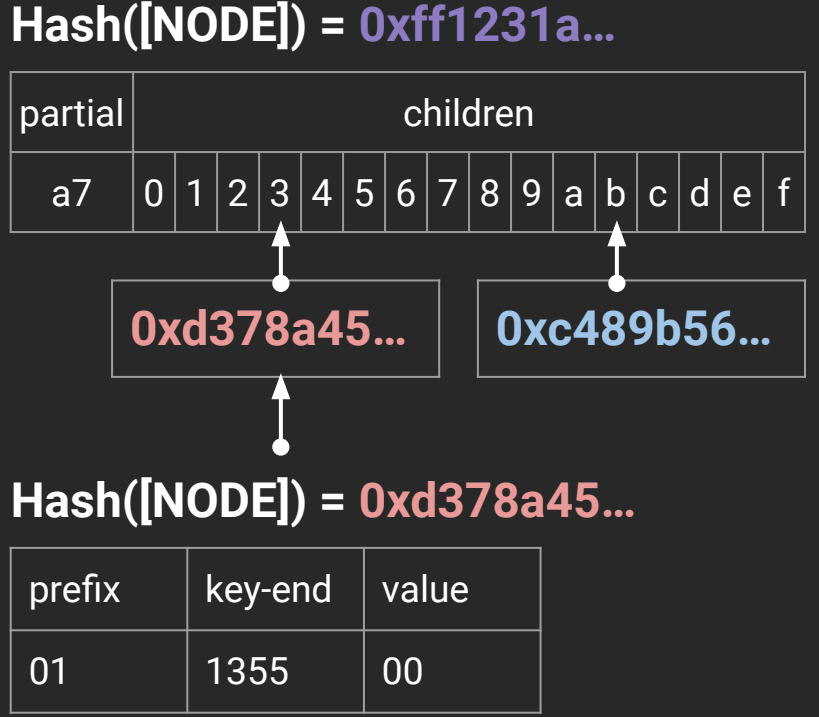
What you just saw

- Patricia provides the trie path.



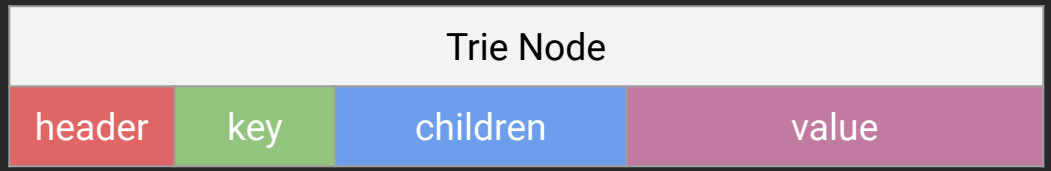
What you just saw

- Patricia provides the trie path.
- Merkle provides the recursive hashing of children nodes into the parent.



Two Kinds of Keys!

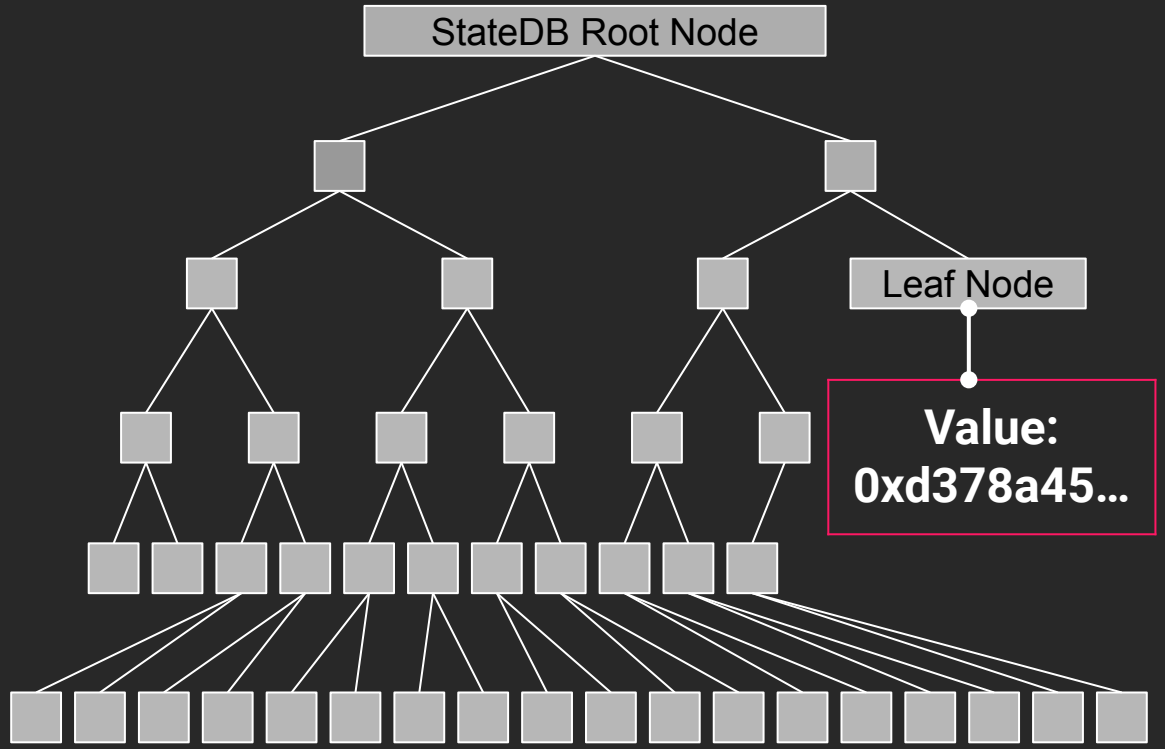
1. Trie key path is set by you! (e.g. “:CODE”)
    - Arbitrary length!
    - Trie Node
        - Header Info
        - Key Info
        - Possible Children
        - Possible Value

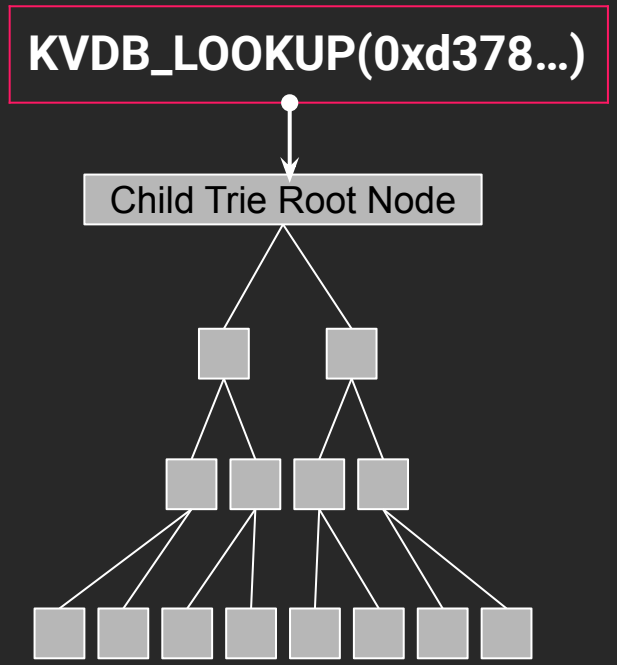


2. KVDB key = Hash([Trie Node])

But wait… there’s more.

Child Trie





* Child tries can be a different format than the Substrate StateDB.

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Prefix Trie

Trie Key Path							Value
a	7						[BRANCH]
a	7	1	1	3	5	5	[LEAF 00]
a	7	7	d	3			[BRANCH w/ VAL 01]
a	7	7	d	3	3	7	[LEAF 02]
a	7	7	d	3	9	7	[LEAF 03]
a	7	f	9	3	6	5	[LEAF 04]

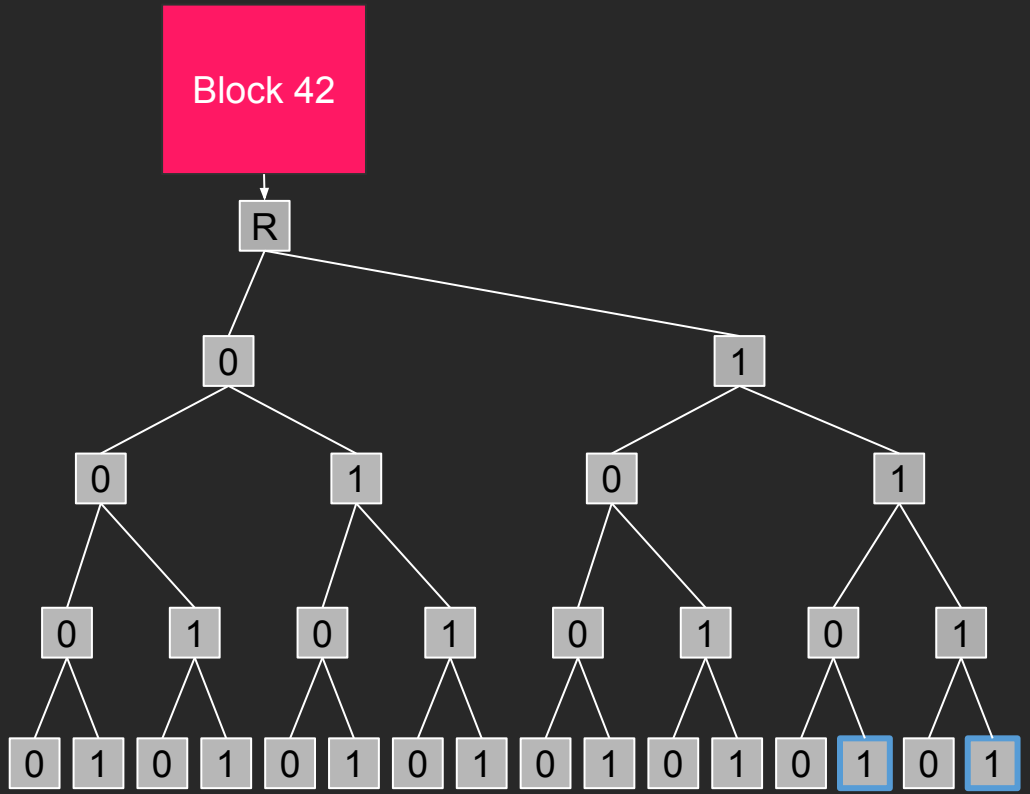
- Similar to Child Trie, but you cannot get the Root Hash.
- Probably something temporary while we fix pruning issues with child trie.

Runtime Storage Trie Path (NEW)

All modules use a prefix trie now! (Long term, they probably become a child trie.)

- Storage Value
  - `twox128(module) + twox128(storagename)`
- linked_map and map
  - `twox128(module) + twox128(storagename) + hasher(key)`
- linked_map head
  - `twox128(module) + twox128("HeadOf" + storagename)`
- double_map
  - `twox128(module) + twox128(storagename) + hasher(key1) + hasher(key2)`

Pruning



- For holding older block states, and cleaning it up.
- Let’s update two values in this trie.

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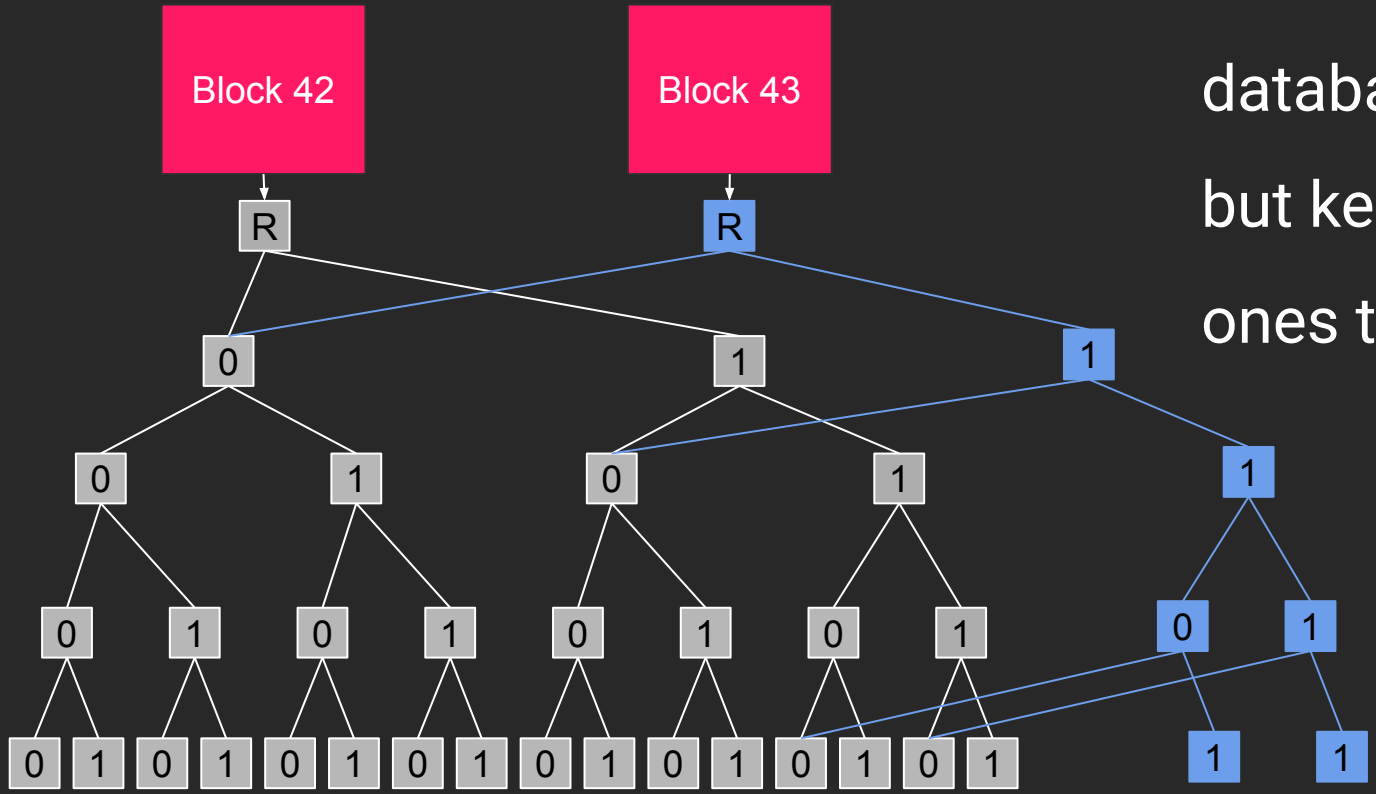
Pruning



We create new database entries, but keep the old ones too!

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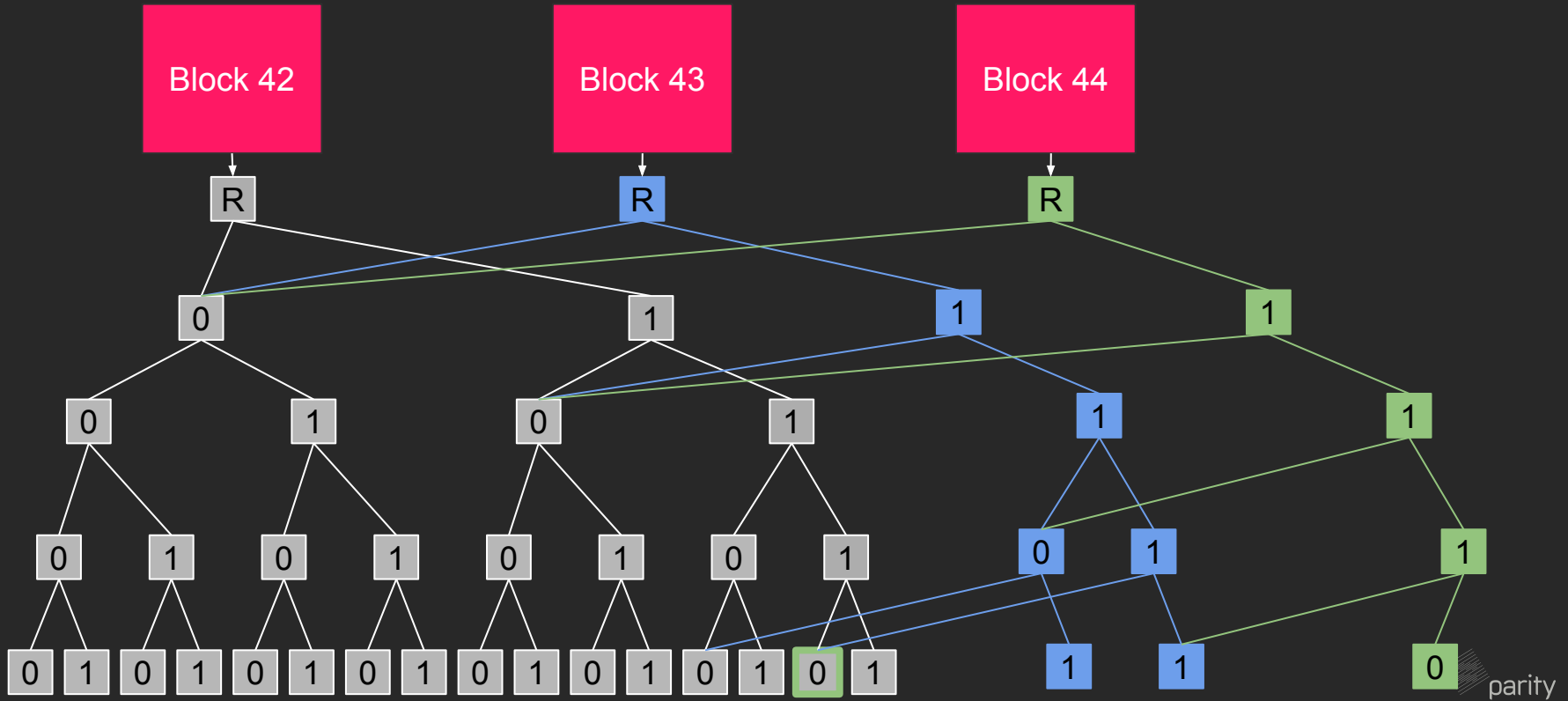
Pruning


Block 42
Block 43
Block 44
R
R
R
0
1
1
1
0
1
0
1
1
1
0
1
0
1
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1
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1
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1
0
1
1
1
0
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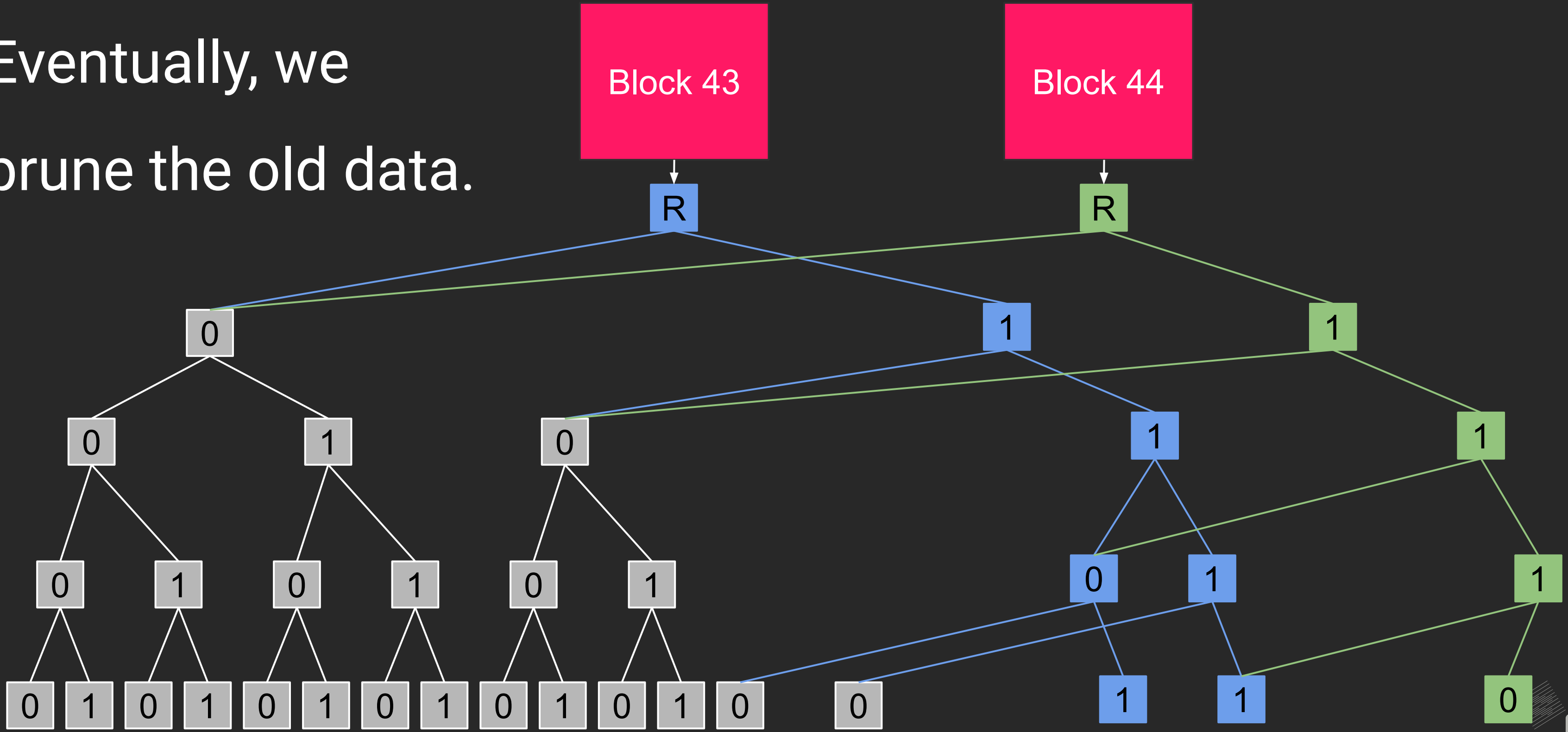
Pruning

Eventually, we

prune the old data.



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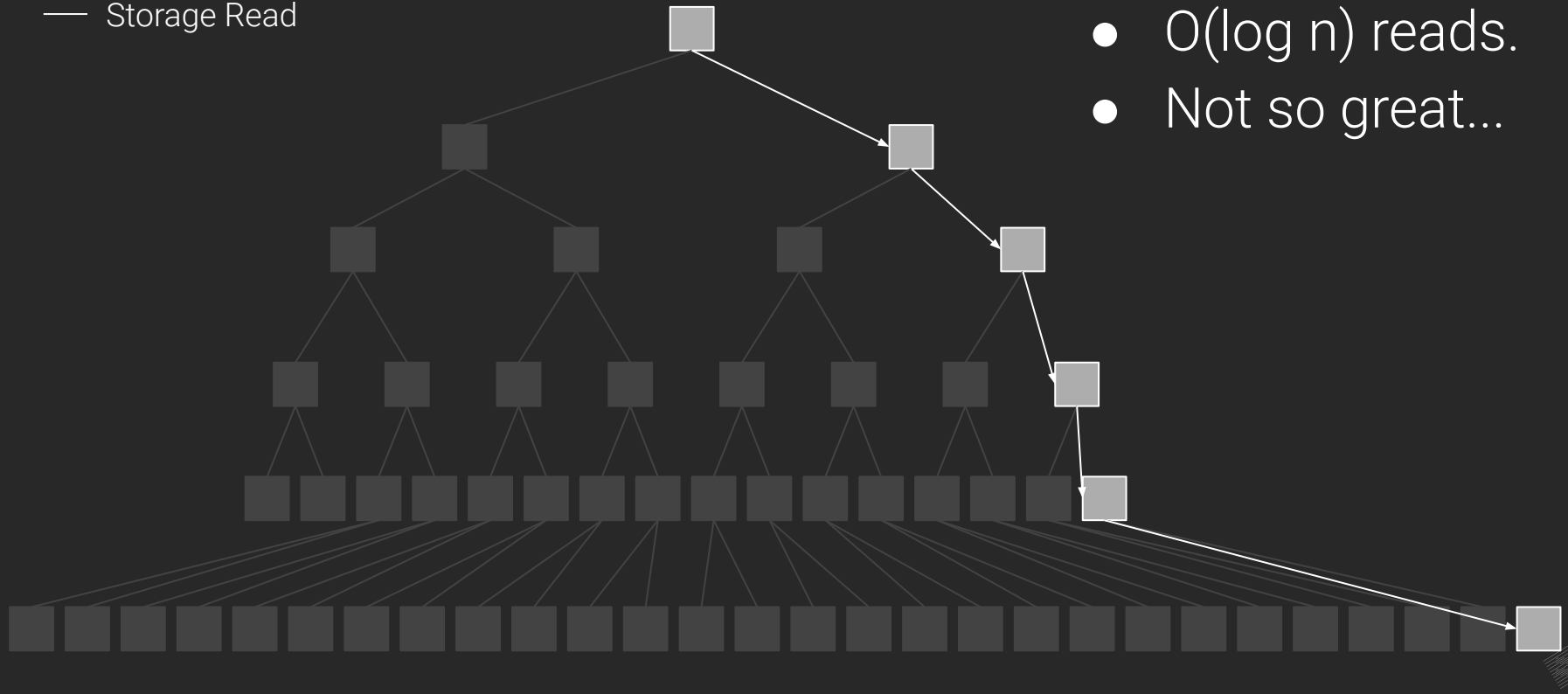
Merkle Trie Complexity

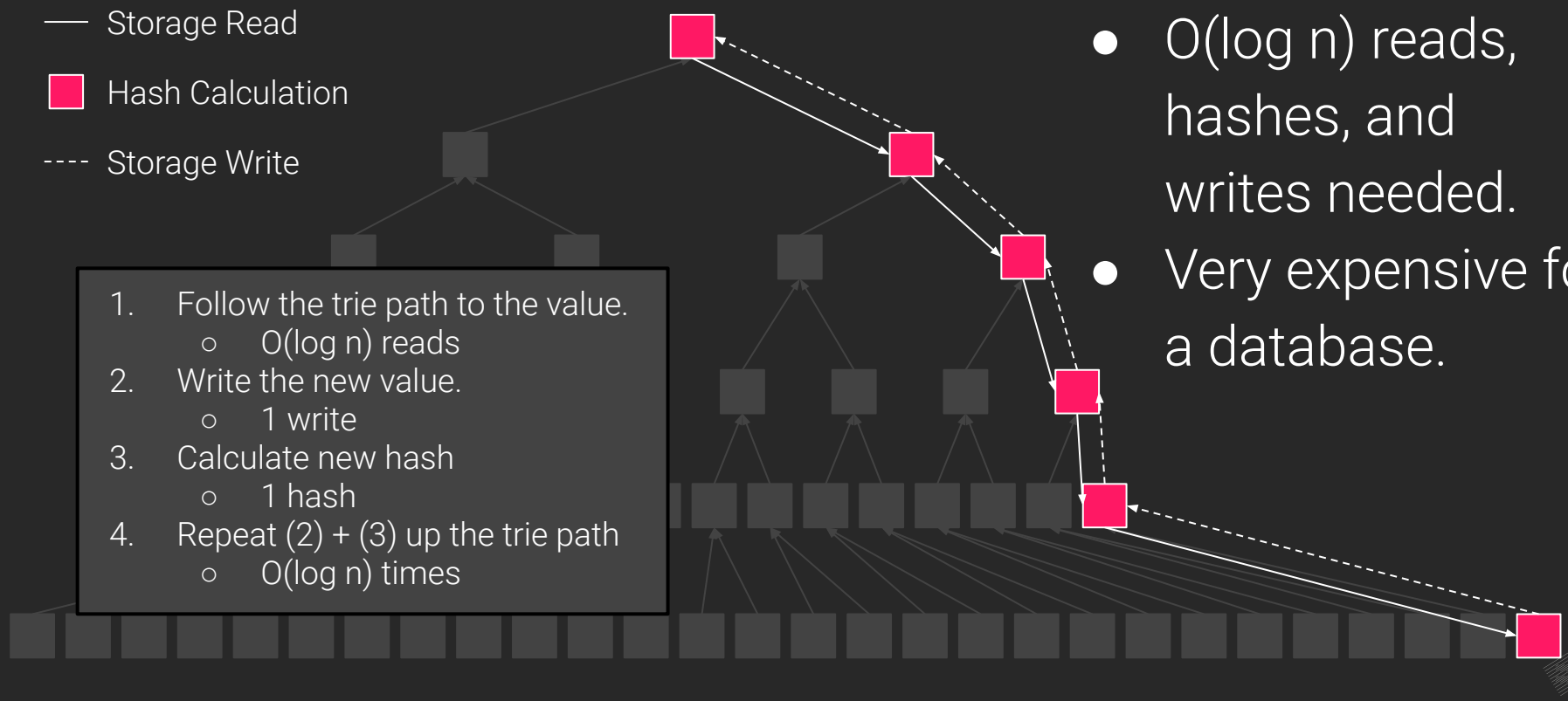
Reading Data

Storage Read



- O(log n) reads.
- Not so great...

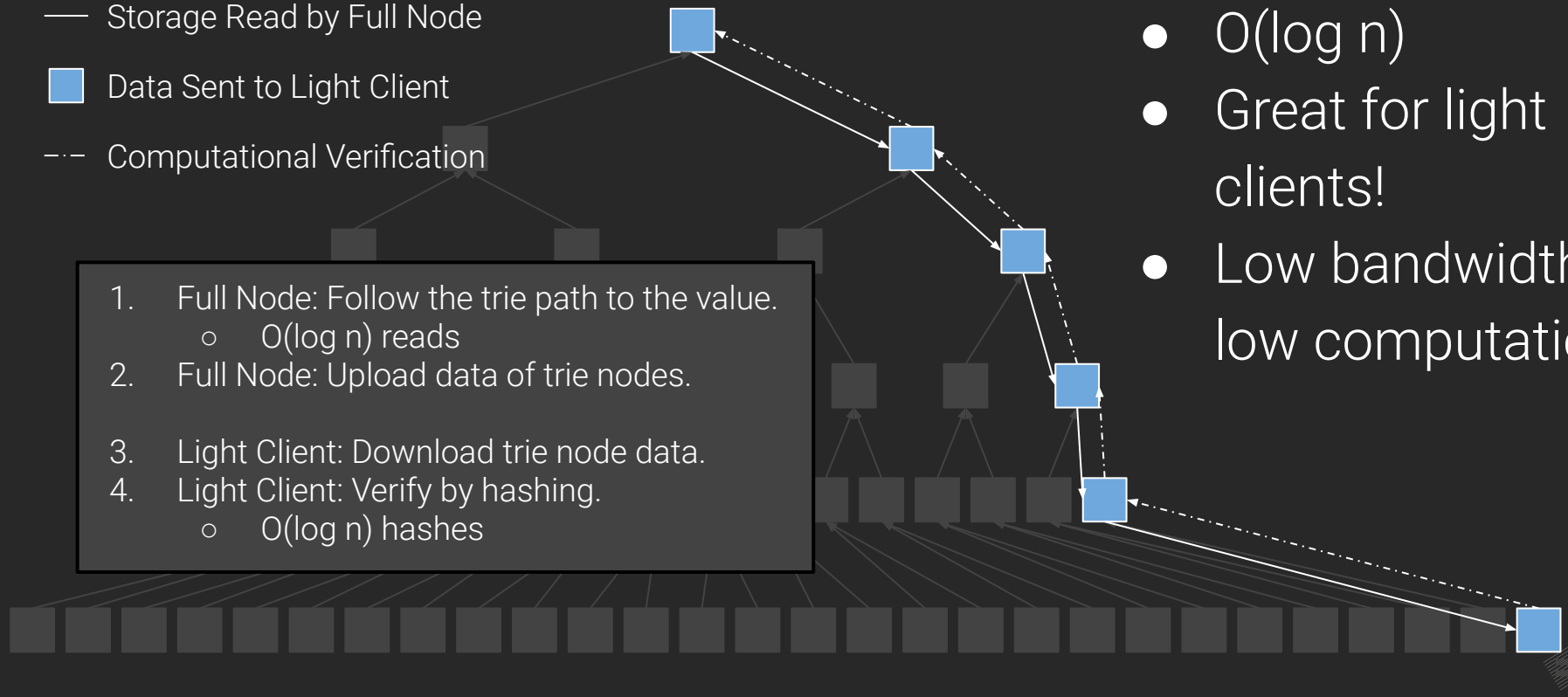
Writing Data



1. Follow the trie path to the value.
   - O(log n) reads
2. Write the new value.
   - 1 write
3. Calculate new hash
   - 1 hash
4. Repeat (2) + (3) up the trie path
   - O(log n) times

- O(log n) reads, hashes, and writes needed.
- Very expensive for a database.

Merkle Proof



1. Full Node: Follow the trie path to the value.
   - O(log n) reads
2. Full Node: Upload data of trie nodes.

3. Light Client: Download trie node data.
4. Light Client: Verify by hashing.
   - O(log n) hashes

- O(log n)
- Great for light clients!
- Low bandwidth, low computation

Best Practices

In general...

Your fundamental goal is to minimize the amount of storage your runtime uses.

You should only store consensus critical data in your runtime storage.

Scenario: Decentralized Blog

- Runtime should be able to come to consensus about the content in a blog post…

★ Store the text on IPFS

★ Store the IPFS hash

➢ DO NOT store the text of the post in the storage!

Struct or Multiple Values?

- Direct costs
  - O(log n) reads to get a value
  - O(log n) writes to update a value
- Indirect costs
  - Increase number of nodes (n)
  - Size of the value

In general… store a struct:

★ Less reads/writes to update multiple values.
★ Less overall nodes in the trie.
★ Adding small items into large items accessed at the same time is essentially free!

➢ Less efficient for single value access.
➢ Upgrades requires storage migration.

Define Your Storage Trie Path Generation

Foo: double_map hasher($hash1) u32, $hash2(u32) => u32

You can control the hashing algorithm used.
By default, these are configured to use Blake2 256.

Final Trie Path:

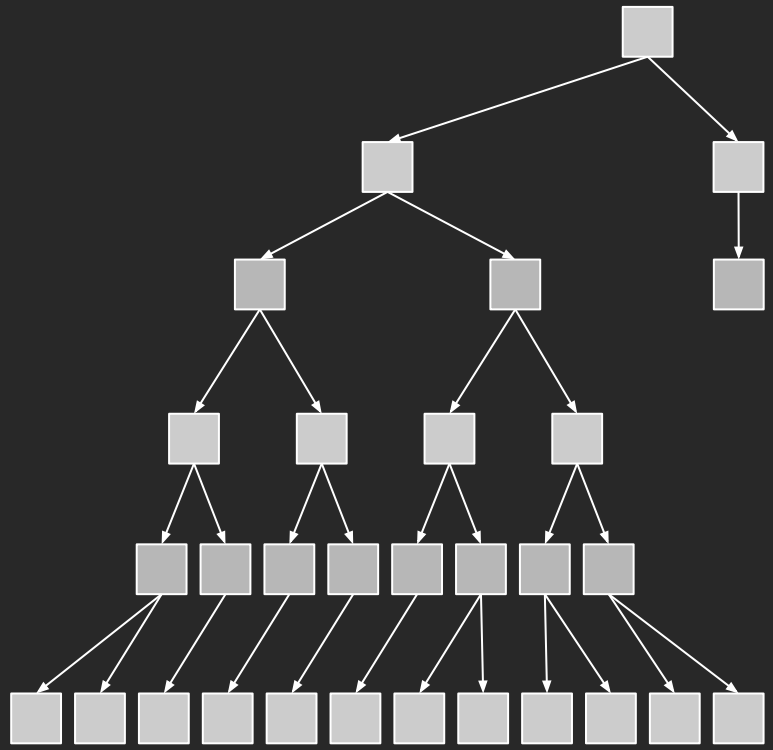
twox128(module) + twox128(storagename) + hasher(key1) + hasher(key2)

XXHash vs Blake2

- What hashing algorithm should I use for trie path generation?

- Blake2
  - Cryptographic but slow…
  - Use when user can influence the input to the hash.
- XXHash (twox)
  - Non-cryptographic, but blazing fast…
  - When you (the runtime developer) controls this value, this is fine!

Unbalanced Trie

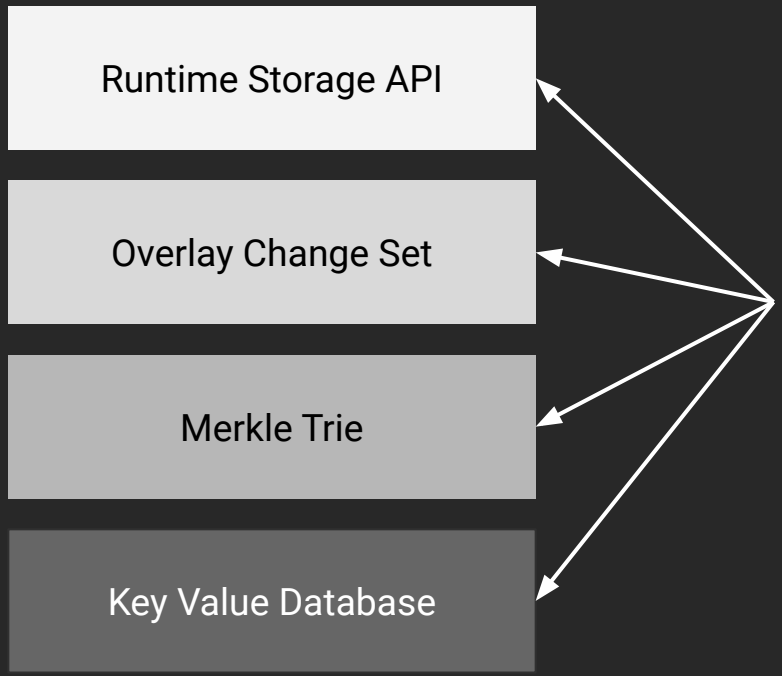


- Can happen if a user can influence the trie path.
- Operations are no longer O(log n)!

Lists

- Vec: For storing a bounded number of values.
  - Good for when you need to change multiple values at a time (single read/write).
  - Enables iteration. Ex: The current validator set.
- Map: For storing an unbounded number of values.
  - Good for random access to data. Ex: User balances.
- Linked Map: For storing unbounded amount of data, but UI or an offchain worker needs to iterate on all the entries.
  - Ex: The list of nominators and their nominations.

Abstractions of Substrate Storage



Think about all the layers when you are writing to Substrate storage.

Questions?

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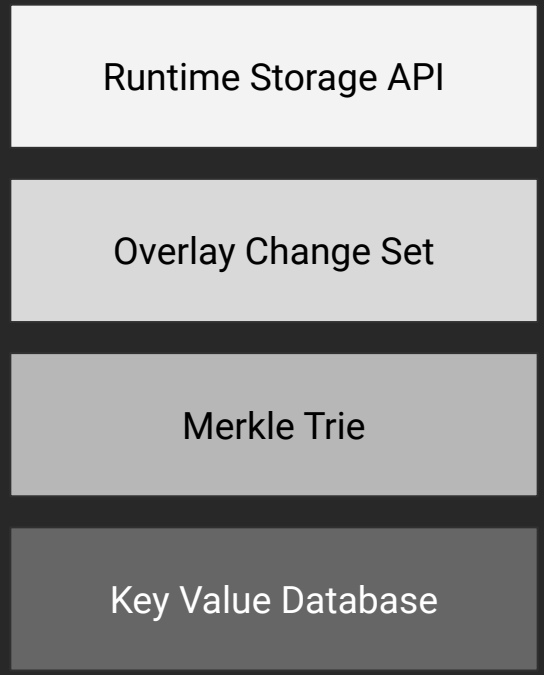
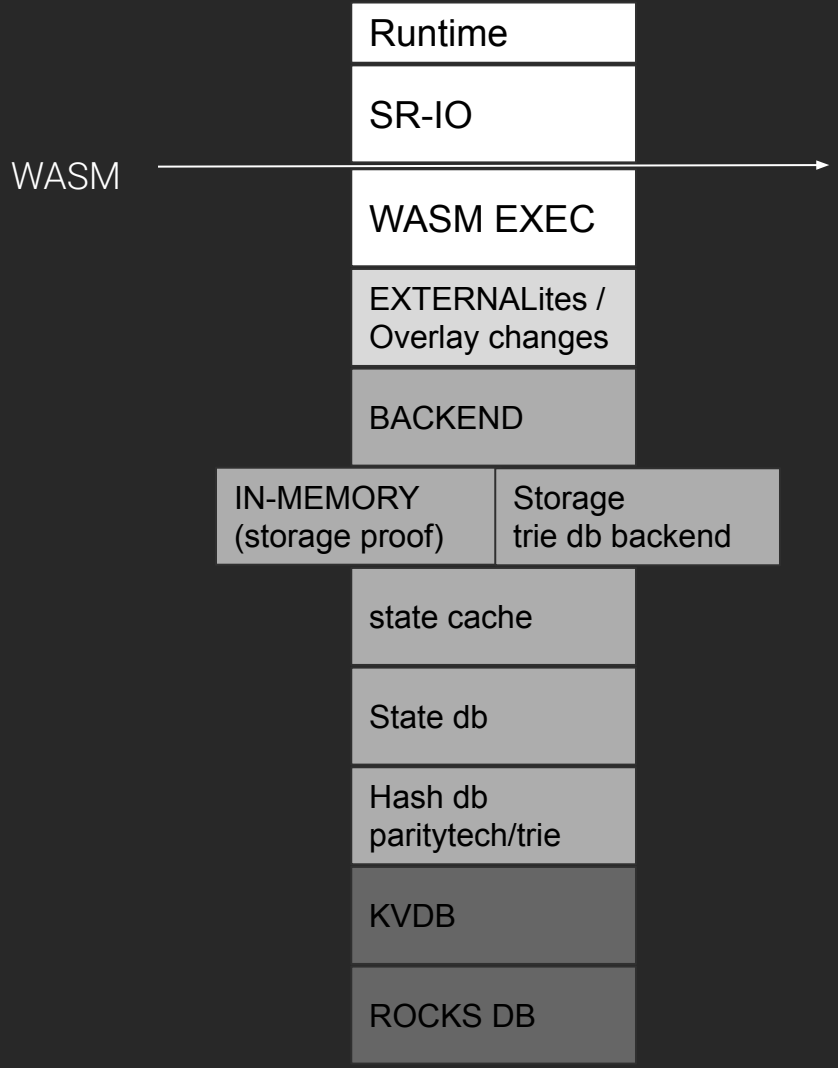
@shawntabrizi


Runtime Storage API
Overlay Change Set
Merkle Trie
Key Value Database


Runtime
SR-IO
WASM
WASM EXEC
EXTERNALites / Overlay changes
BACKEND
IN-MEMORY (storage proof)
Storage trie db backend
state cache
State db
Hash db paritytech/trie
KVDB
ROCKS DB

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